

**Clean version of the Claims**

*Clean version of the claims with all of the changes to be made vis-à-vis the U.S. Patent 5,848,159, as follows:*

1. (Twice Amended) A method for communications of a message cryptographically processed with RSA (Rivest, Shamir & Adleman) public key encryption, comprising the steps of: developing  $k$  distinct random prime numbers  $p_1, p_2, \dots, p_k$ , where  $k$  is an integer greater than 2; providing a number  $e$  relatively prime to  $(p_1 - 1) \cdot (p_2 - 1) \cdot \dots \cdot (p_k - 1)$ ; providing a composite number  $n$  equaling the product  $p_1 \cdot p_2 \cdot \dots \cdot p_k$ ; receiving a ciphertext word signal  $C$  which is formed by encoding a plaintext message word signal  $M$  to a ciphertext word signal  $C$ , where  $M$  corresponds to a number representative of the message and

$$0 \leq M \leq n-1,$$

where  $C$  is a number representative of an encoded form of the plaintext message word signal  $M$  such that

$C \equiv M^e \pmod{n}$ , and where  $e$  is associated with an intended recipient of the ciphertext word signal  $C$ ; and

deciphering the received ciphertext word signal  $C$  at the intended recipient having available to it the  $k$  distinct random prime numbers  $p_1, p_2, \dots, p_k$ .

2. (Twice Amended) The method according to claim 1, wherein the deciphering step includes establishing a number,  $d$ , as a multiplicative inverse of  $e \pmod{\text{lcm}((p_1 - 1), (p_2 - 1), \dots, (p_k - 1)))}$ , and decoding the ciphertext word signal  $C$  to the plaintext message word signal  $M$  where  $M \equiv C^d \pmod{n}$ .

3. (Twice Amended) A method for communications of a message signal  $M_i$  cryptographically processed with RSA public key encryption in a system having  $j$  terminals, each terminal being characterized by an encoding key  $E_i = (e_i, n_i)$  and a decoding key  $D_i = (d_i, n_i)$ , where  $i = 1, 2, \dots, j$ ,

and the message signal  $M_i$  corresponds to a number representative of a message-to-be-received from the  $i^{\text{th}}$  terminal, the method comprising the steps of:  
establishing  $n_i$  where  $n_i$  is a composite number of the form

$$n_i = p_{i,1} \cdot p_{i,2} \cdot \dots \cdot p_{i,k}$$

where  $k$  is an integer greater than 2,

$p_{i,1}, p_{i,2}, \dots, p_{i,k}$  are distinct random prime numbers,

$e_i$  is relatively prime to  $\text{lcm}(p_{i,1} - 1, p_{i,2} - 1, \dots, p_{i,k} - 1)$ , and

$d_i$  is selected from the group consisting of the class of numbers equivalent to a multiplicative inverse of

$$e_i \pmod{(\text{lcm}((p_{i,1} - 1), (p_{i,2} - 1), \dots, (p_{i,k} - 1)))};$$

receiving by a recipient terminal ( $i = y$ ) from a sender terminal ( $i = x, x \neq y$ ) a ciphertext signal  $C_x$  formed by encoding a digital message word signal  $M_x$ , wherein the encoding includes

transforming said message word signal  $M_x$  to one or more message block word signals  $M_x''$ , each block word signal  $M_x''$  corresponding to a number representative of a portion of said message word signal  $M_x$  in the range  $0 \leq M_x'' \leq n_y - 1$ , and

transforming each of said message block word signals  $M_x''$  to a ciphertext word signal  $C_x$  that corresponds to a number representative of an encoded form of said message block word signal  $M_x''$  where

$$C_x \equiv M_x''^{e_y} \pmod{n_y}; \text{ and}$$

deciphering the received ciphertext word signal  $C_x$  at the recipient terminal having available to it the  $k$  distinct random prime numbers  $p_{y,1}, p_{y,2}, \dots, p_{y,k}$  for establishing its  $d_y$ .

4. (Twice Amended) A system for communications of a message cryptographically processed with an RSA public key encryption, comprising:

a communication channel for transmitting a ciphertext word signal  $C$ ;

encoding means coupled to said channel and adapted for transforming a transmit message word signal  $M$  to the ciphertext word signal  $C$  using a composite number,  $n$ , where  $n$  is a product of the form

$$n = p_1 \cdot p_2 \cdot \dots \cdot p_k$$

k is an integer greater than 2, and

$p_1, p_2, \dots, p_k$  are distinct random prime numbers,

where the transmit message word signal M corresponds to a number representative of the message and

$0 \leq M \leq n-1$

where the ciphertext word signal C corresponds to a number representative of an encoded form of said message through a relationship of the form

$C \equiv M^e \pmod{n}$ , and

where e is a number relatively prime to  $\text{lcm}(p_1 - 1, p_2 - 1, \dots, p_k - 1)$ ; and

decoding means coupled to said channel and adapted for receiving the ciphertext word signal C from said channel and, having available to it the k distinct random prime numbers  $p_1, p_2, \dots, p_k$ , for transforming the ciphertext word signal C to a receive message word signal M' where M' corresponds to a number representative of a decoded form of the ciphertext word signal C through a relationship of the form

$M \equiv C^d \pmod{n}$

where d is selected from the group consisting of a class of numbers equivalent to a multiplicative inverse of

$e(\text{mod}(\text{lcm}((p_1 - 1), (p_2 - 1), \dots, (p_k - 1))))$ .

5. (Twice Amended) A system for communications of a message cryptographically processed with an RSA public key encryption, the system having a plurality of terminals coupled by a communications channel, comprising:

a first terminal of the plurality of terminals characterized by an encoding key

$E_A = (e_A, n_A)$  and a decoding key  $D_A = (d_A, n_A)$ ,

where  $n_A$  is a composite number of the form

$n_A = p_{A,1} \cdot p_{A,2} \cdot \dots \cdot p_{A,k}$

where

k is an integer greater than 2,

$p_{A,1}, p_{A,2}, \dots, p_{A,k}$  are distinct random prime numbers,

$e_A$  is relatively prime to

$\text{lcm}(p_{A,1} - 1, p_{A,2} - 1, \dots, p_{A,k} - 1)$ , and

$d_A$  is selected from the group consisting of the class of numbers equivalent to a multiplicative inverse of  $e_A \pmod{\text{lcm}((p_{A,1}-1), (p_{A,2}-1), \dots, (p_{A,k}-1))}$ ; and

a second terminal of the plurality of terminals having

blocking means for transforming a first message, which is to be transmitted on said communications channel from said second terminal to said first terminal, into one or more transmit message word signals  $M_B$ , where each  $M_B$  corresponds to a number representative of said first message in the range

$$0 \leq M_B \leq n_A - 1,$$

encoding means coupled to said channel and adapted for transforming each transmit message word signal  $M_B$  to a ciphertext word signal  $C_B$  that corresponds to a number representative of an encoded form of said first message through a relationship of the form

$$C_B \equiv M_B^{e_A} \pmod{n_A},$$

said first terminal having

decoding means coupled to said channel and adapted for receiving each of said ciphertext word signals  $C_B$  from said channel and, having available to it the  $k$  distinct random prime numbers  $p_{A,1}, p_{A,2}, \dots, p_{A,k}$ , for transforming each of said ciphertext word signals  $C_B$  to a receive message word signal  $M'_B$ , and means for transforming said receive message word signal  $M'_B$  to said first message, where  $M'_B$  corresponds to a number representative of a decoded form of  $C_B$  through a relationship of the form

$$M'_B \equiv C_B^{d_A} \pmod{n_A}.$$

6. (Twice Amended) The system according to claim 5 wherein said second terminal is characterized by an encoding key  $E_B = (e_B, n_B)$  and a decoding key  $D_B = (d_B, n_B)$ , where  $n_B$  is a composite number of the form

$$n_B = p_{B,1} \cdot p_{B,2} \cdot \dots \cdot p_{B,k}$$

where  $k$  is an integer greater than 2,

$p_{B,1}, p_{B,2}, \dots, p_{B,k}$  are distinct random prime numbers,

$e_B$  is relatively prime to

$\text{lcm}(p_{B,1}-1, p_{B,2}-1, \dots, p_{B,k}-1)$ , and

$d_B$  is selected from the group consisting of a class of numbers equivalent to a multiplicative inverse of

$e_B \pmod{\text{lcm}((p_{B,1}-1), (p_{B,2}-1), \dots, (p_{B,k}-1))}$ ,

said first terminal further having

blocking means for transforming a second message, [-to-be-transmitted] which is to be transmitted on said communications channel from said first terminal to said second terminal, to one or more transmit message word signals  $M_A$ , where each  $M_A$  corresponds to a number representative of said message in the range

$$0 \leq M_A \leq n_B - 1$$

encoding means coupled to said channel and adapted for transforming each transmit message word signal  $M_A$  to a ciphertext word signal  $C_A$  and for transmitting  $C_A$  on said channel, where  $C_A$  corresponds to a number representative of an encoded form of said second message through a relationship of the form

$$C_A \equiv M_A^{e_B} \pmod{n_B}$$

said second terminal further having

decoding means coupled to said channel and adapted for receiving each of said ciphertext word signals  $C_A$  from said channel and, having available to it the  $k$  distinct random prime numbers  $p_{B,1}, p_{B,2}, \dots, p_{B,k}$ , for transforming each of said ciphertext word signals to a receive message word signal  $M'_A$ , and means for transforming said receive message word signals  $M'_A$  to said second message, where  $M'_A$  corresponds to a number representative of a decoded form of  $C_A$  through a relationship of the form

$$M'_A \equiv C_A^{d_B} \pmod{n_B}.$$

7. (Amended) A method for communications of a message cryptographically processed with an RSA public key encryption, comprising the steps of:

developing  $k$  factors of a composite number  $n$ , the  $k$  factors being distinct random prime numbers and  $k$  is an integer larger than two ( $k > 2$ );

providing a number  $e$  relatively prime to a lowest common multiplier of the  $k$  factors;

providing the composite number  $n$ ;

receiving a ciphertext word signal C which is formed by encoding a digital message word signal M to the ciphertext word signal C, where said digital message word signal M corresponds to a number representative of said message and

$$0 \leq M \leq n-1,$$

where said ciphertext word signal C corresponds to a number representative of an encoded form of said message through a relationship of the form

$$C \equiv a_e M^e + a_{e-1} M^{e-1} + \dots + a_0 \pmod{n}$$

where e and  $a_e, a_{e-1}, \dots, a_0$  are numbers; and

deciphering the received ciphertext word signal C at an intended recipient with knowledge of the k factors.

8. Cancelled.

9. (Twice Amended) A system for communications of message signals cryptographically processed with RSA public key encryption, comprising:

j terminals including first and second terminals, each of the j terminals being characterized by an encoding key  $E_i = (e_i, n_i)$  and decoding key  $D_i = (d_i, n_i)$ , where  $i = 1, 2, \dots, j$ , each of the j terminals being adapted to transmit a particular one of the message signals where an  $i^{\text{th}}$  message signal  $M_i$  is transmitted from an  $i^{\text{th}}$  terminal, and

$$0 \leq M_i \leq n_i - 1,$$

$n_i$  being a composite number of the form

$$n_i = p_{i,1} \cdot p_{i,2} \cdot \dots \cdot p_{i,k}$$

where

$k$  is an integer greater than 2,

$p_{i,1}, p_{i,2}, \dots, p_{i,k}$  are distinct random prime numbers,

$e_i$  is relatively prime to

$$\text{lcm}(p_{i,1}-1, p_{i,2}-1, \dots, p_{i,k}-1), \text{ and}$$

$d_i$  is selected from the group consisting of the class of numbers equivalent to a multiplicative inverse of

$$e_i \pmod{\text{lcm}((p_{i,1}-1), (p_{i,2}-1), \dots, (p_{i,k}-1))};$$

said first terminal including

means for encoding a digital message word signal  $M_1$  to be transmitted from said first terminal ( $i=1$ ) to said second terminal ( $i=2$ ), said encoding means transforming said digital message word signal  $M_1$  to a signed message word signal  $M_{1s}$  using a relationship of the form

$$M_{1s} \equiv M_1^{d_1} \pmod{n_1}; \text{ and}$$

means for transmitting said signed message word signal  $M_{1s}$  from said first terminal to said second terminal, wherein said second terminal includes

means for decoding said signed message word signal  $M_{1s}$  to said digital message word signal  $M_1$ .

10. (Twice Amended) The system of claim 9, wherein the means for decoding said signed message word signal  $M_{As}$  includes means for transforming said signed message word signal  $M_{As}$  using a relationship of the form

$$M_1 \equiv M_{1s}^{e_1} \pmod{n_1}.$$

11. (Twice Amended) A communications system for transferring a message signal cryptographically processed with RSA public key encryption, the communications system comprising:

$j$  communication stations including first and second stations, each of the  $j$  communication stations being characterized by an encoding key  $E_i = (e_i, n_i)$  and a decoding key  $D_i = (d_i, n_i)$ , where  $i=1, 2, \dots, j$ , each of the  $j$  communication stations being adapted to transmit a particular one of the message signals where an  $i^{\text{th}}$  message signal  $M_i$  is received from an  $i^{\text{th}}$  communication station, and

$$0 \leq M_i \leq n_i - 1$$

$n_i$  being a composite number of the form

$$n_i = p_{i,1} \cdot p_{i,2} \cdot \dots \cdot p_{i,k}$$

where

$k$  is an integer greater than 2,

$p_{i,1}, p_{i,2}, \dots, p_{i,k}$  are distinct random prime numbers,

$e_i$  is relatively prime to  $\text{lcm}(p_{i,1} - 1, p_{i,2} - 1, \dots, p_{i,k} - 1)$ , and

$d_i$  is selected from the group consisting of the class of numbers equivalent to a multiplicative inverse of

$$e_i \pmod{\text{lcm}((p_{i,1}-1), (p_{i,2}-1), \dots, (p_{i,k}-1))},$$

said first station including

means for encoding a digital message word signal  $M_1$  to be transmitted from said first station ( $i=1$ ) to said second station ( $i=2$ ),

means for transforming said digital message word signal  $[M_A] M_1$  to one or more message block word signals  $[M_A'] M_1''$ , each block word signal  $[M_A'] M_1''$  being a number representative of a portion of said message word signal  $M_1$  in the range

$$0 \leq M_1'' \leq n_2-1, \text{ and}$$

means for transforming each of said message block word signals  $M_1''$  to a ciphertext word signal  $C_1$  using a relationship of the form

$$C_1 \equiv M_1''^{e_2} \pmod{n_2}; \text{ and}$$

means for transmitting said ciphertext word signals  $C_1$  from said first station to said second station, wherein said second station includes

means for deciphering said ciphertext word signals  $C_1$  using  $p_{2,1}, p_{2,2}, \dots, p_{2,k}$  to produce said message word signal  $M_1$ .

12. (Twice Amended) The communications system of claim 11, wherein the deciphering means includes

means for decoding said ciphertext word signals  $C_1$  to said message block word signals  $M_1''$  using a relationship of the form

$$M_1'' \equiv C_1^{d_2} \pmod{n_2}, \text{ and}$$

means for transforming said message block word signals  $M_1''$  to said message word signal  $M_1$ .

13. (Twice Amended) A system for communications of a message cryptographically processed with RSA public key encryption, comprising:

a first station; and

a second station communicatively connected to the first station,

the first station having

encoding means for transforming a transmit message word signal  $M$  to a ciphertext word signal  $C$  where the transmit message word signal  $M$  corresponds to a number representative of a message and

$$0 \leq M \leq n-1$$

$n$  being a composite number formed as a product of at least 3 factors, the at least 3 factors being distinct random prime numbers, and where the ciphertext word signal  $C$  corresponds to a number representative of an encoded form of said message through a relationship of the form

$$C \equiv a_e M^e + a_{e-1} M^{e-1} + \dots + a_0 \pmod{n}$$

where  $e$  and  $a_e, a_{e-1}, \dots, a_0$  are numbers; and

means for transmitting the ciphertext word signal  $C$  to the second station, wherein the second station includes means for deciphering the ciphertext word signal  $C$  using the at least 3 factors to produce the message.

New Claims:

14. (Amended) A method of communicating a message cryptographically processed with an RSA public key encryption, comprising the steps of:

selecting a public key portion  $e$  associated with a recipient intended for receiving the message;

developing  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ , where  $k \geq 3$ , and checking that each of the  $k$  distinct random prime numbers minus 1,  $p_1-1, p_2-1, \dots, p_k-1$ , is relatively prime to the public key portion  $e$ ;

computing a composite number,  $n$ , as a product of the  $k$  distinct random prime numbers;

receiving a ciphertext message formed by encoding a plaintext message data  $M$  to the ciphertext message data  $C$  using a relationship of the form  $C \equiv M^e \pmod{n}$ , where  $M$  represents the

message, where  $0 \leq M \leq n-1$  and where the sender knows  $n$  and the public key portion  $e$  but has no access to the  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ ; and deciphering at the recipient the received ciphertext message data  $C$  to produce the message, the recipient having access to the  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ .

15. (Amended) The method according to claim 14, comprising the further step of:

establishing a private key portion  $d$  by a relationship to the public key portion  $e$  in the form of

$$d \equiv e^{-1} \pmod{((p_1-1) \cdot (p_2-1) \cdots (p_k-1))},$$

wherein the deciphering step includes decoding the ciphertext message data  $C$  to the plaintext message data  $M$  using a relationship of the form  $M \equiv C^d \pmod{n}$ .

16. (Amended) A method of communicating a message cryptographically processed with RSA public key encryption, comprising the steps of:

selecting a public key portion  $e$ ;

developing  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ , where  $k \geq 3$ , and checking that each of the  $k$  distinct random prime numbers minus 1,  $p_1-1, p_2-1, \dots, p_k-1$ , is relatively prime to the public key portion  $e$ ;

establishing a private key portion  $d$  by a relationship to the public key portion  $e$  in the form of

$$d \equiv e^{-1} \pmod{((p_1-1) \cdot (p_2-1) \cdots (p_k-1))};$$

computing a composite number,  $n$ , as a product of the  $k$  distinct random prime numbers;

receiving a ciphertext message data  $C$  representing an encoded form of a plaintext message data  $M$ ; and

decoding the received ciphertext message data  $C$  to the plaintext message data  $M$  using a relationship of the form  $M \equiv C^d \pmod{n}$ , the decoding performed by a recipient owning the private key portion  $d$  and having access to the  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ .

17. (Amended) The method according to claim 16, wherein the ciphertext message data  $C$  is formed by encoding the plaintext message data  $M$  to the ciphertext message data  $C$  using a relationship of the form  $C \equiv M^e \pmod{n}$ , wherein  $0 \leq M \leq n-1$  and wherein  $n$  and the public key portion  $e$  are accessible to the sender although it has no access to the  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ .

18. (Amended) A method of communicating a message cryptographically processed with RSA public key encryption, comprising the steps of:

selecting a public key portion  $e$ ;

developing  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ , where  $k \geq 3$ , and checking that each of the  $k$  distinct random prime numbers minus 1,  $p_1-1, p_2-1, \dots, p_k-1$ , is relatively prime to the public key portion  $e$ ;

establishing a private key portion  $d$  by a relationship to the public key portion  $e$  of the form

$$d \equiv e^{-1} \pmod{(p_1-1) \cdot (p_2-1) \cdots (p_k-1)};$$

computing a composite number,  $n$ , as a product of the  $k$  distinct random prime numbers;

encoding a plaintext message data  $M$  with the private key portion  $d$  to produce a signed message  $M_s$  using a relationship of the form  $M_s \equiv M^d \pmod{n}$ , where  $0 \leq M \leq n-1$

receiving the signed message  $M_s$ ; and

deciphering the signed message to produce the plaintext message data  $M$ .

19. (Amended) The method of claim 18, wherein the deciphering step includes:

decoding the signed message  $M_s$  with the public key portion  $e$  to produce the plaintext message data  $M$  using a relationship of the form  $M \equiv M_s^e \pmod{n}$ .

20. (Amended) A method for communicating a message cryptographically processed with RSA public key encryption, comprising the steps of:

sending to a recipient a cryptographically processed message formed by

assigning a number  $M$  to represent the message in plaintext message form, and

cryptographically transforming the assigned number  $M$  from the plaintext message form

to a number  $C$  that represents the message in an encoded form, wherein the number  $C$  is a function of

the assigned number  $M$ ,

a number  $n$  that is a composite number equaling the product of at least three distinct random prime numbers, wherein  $0 \leq M \leq n-1$ , and

an exponent  $e$  that is a number relatively prime to a lowest common multiplier of the at least three distinct random prime numbers,

wherein the number  $n$  and exponent  $e$  having been obtained by the sender are associated with the recipient to which the message is intended; and

receiving the cryptographically processed message which is decipherable by the recipient based on

the number  $n$ ,

another exponent  $d$ , and

the number  $C$ ,

wherein the exponent  $d$  is a function of the exponent  $e$  and the at least three distinct random prime numbers.

21. (Amended) The method according to claim 20,

wherein the cryptographically transforming step includes using a relationship of the form  $C \equiv M^e \pmod{n}$ ,

wherein the exponent  $d$  is established based on the at least three distinct random prime numbers,  $p_1, p_2, \dots, p_k$  using a relationship of the form  $d \equiv e^{-1} \pmod{((p_1-1) \cdot (p_2-1) \cdots (p_k-1))}$ ,

and

wherein the cryptographically processed message is deciphered using a relationship of the form  
 $M \equiv C^d \pmod{n}$ .

22. (Amended) A method for communicating a message cryptographically processed with RSA public key encryption, comprising the steps of:

receiving from a sender a cryptographically processed message, in the form of a number  $C$ , which is decipherable by the recipient based on a number  $n$ , an exponent  $d$ , and the number  $C$ ; and

deciphering the cryptographically processed message,

wherein a number  $M$  represents a plaintext form of the message, wherein the number  $C$  represents a cryptographically encoded form of the message and is a function of the number  $M$ ,

the number  $n$  that is a composite number equaling the product of at least three distinct random prime numbers, wherein  $0 \leq M \leq n-1$ , and

an exponent  $e$  that is a number relatively prime to a lowest common multiplier of the at least three distinct random prime numbers,

wherein the number  $n$  and exponent  $e$  are associated with the recipient to which the message is intended, and

wherein the exponent  $d$  is a function of the exponent  $e$  and the at least three distinct random prime numbers.

23. (Amended) The method according to claim 22,

wherein the number  $C$  is formed using a relationship of the form  $C \equiv M^e \pmod{n}$ ,

wherein the exponent  $d$  is established based on the at least three distinct random prime numbers,  $p_1, p_2, \dots, p_k$  using a relationship of the form  $d \equiv e^{-1} \pmod{((p_1-1) \cdot (p_2-1) \cdots (p_k-1))}$ ,

and wherein the number  $M$  is obtained using a relationship of the form  $M \equiv C^d \pmod{n}$ .

24. (Amended) The method according to claim 21,

wherein  $p$  and  $q$  are a pair of prime numbers the product of which equals  $n$ ,  
wherein the deciphering the number  $C$  to derive the number  $M$  is divided into subtasks, one  
subtask for each of the  $k$  distinct random prime numbers,  
wherein the  $k$  distinct random prime numbers are each smaller than  $p$  and  $q$ ,  
whereby for a given length of  $n$  it takes fewer computational cycles to perform the deciphering  
relative to the number of computational cycles for performing such deciphering if the pair  
of prime numbers  $p$  and  $q$  were used instead.

25. (Amended) The method according to claim 22,  
wherein  $p$  and  $q$  are a pair of prime numbers the product of which equals  $n$ ,  
wherein the deciphering the number  $C$  to derive the number  $M$  is divided into subtasks, one  
subtask for each of the  $k$  distinct random prime numbers,  
wherein the  $k$  distinct random prime numbers are each smaller than  $p$  and  $q$ ,  
whereby for a given length of  $n$  it takes fewer computational cycles to perform the deciphering  
relative to the number of computational cycles for performing such deciphering if the pair of  
prime numbers  $p$  and  $q$  were used instead.

26. (Amended) The method according to claim 20,  
wherein  $p$  and  $q$  are a pair of prime numbers the product of which equals  $n$ , and  
wherein developing the at least three distinct random prime numbers and computing  $n$  is  
performed, including for  $n$  that is more than 600 digits long, in less time than it takes to develop  
the pair of prime numbers  $p$  and  $q$  and compute that  $n$ .

27. (Amended) The method according to claim 22,  
wherein  $p$  and  $q$  are a pair of prime numbers the product of which equals  $n$ , and  
wherein developing the at least three distinct random prime numbers and computing  $n$  is  
performed, including for  $n$  that is more than 600 digits long, in less time than it takes to develop  
the pair of prime numbers  $p$  and  $q$  and compute that  $n$ .

28. (Amended) The method according to claim 14,  
wherein  $p$  and  $q$  are a pair of prime numbers the product of which equals  $n$ ,

wherein the deciphering step is divided into sub-steps, one sub-step for each of the  $k$  distinct random prime numbers,

wherein the  $k$  distinct random prime numbers are each smaller than  $p$  and  $q$ ,

whereby for a given length of  $n$  it takes fewer computational cycles to perform the deciphering step relative to the number of computational cycles for performing such deciphering step if the pair of prime numbers  $p$  and  $q$  were used instead.

29. (Amended) The method according to claim 14,

wherein  $p$  and  $q$  are a pair of prime numbers the product of which equals  $n$ , and

wherein developing the  $k$  distinct random prime numbers and computing the composite number  $n$  are performed, including for  $n$  that is more than 600 digits long, in less time than it takes to develop the pair of prime numbers  $p$  and  $q$  and compute that  $n$ .

30. (Amended) The method according to claim 16,

wherein  $p$  and  $q$  are a pair of prime numbers the product of which equals  $n$ ,

wherein the decoding step is divided into sub-steps, one sub-step for each of the  $k$  distinct random prime numbers,

wherein the  $k$  distinct random prime numbers are each smaller than  $p$  and  $q$ ,

whereby for a given length of  $n$  it takes fewer computational cycles to perform the decoding step relative to the number of computational cycles for performing such decoding step if the pair of prime numbers  $p$  and  $q$  were used instead.

31. (Amended) The method according to claim 16,

wherein  $p$  and  $q$  are a pair of prime numbers the product of which equals  $n$ , and

wherein developing the  $k$  distinct random prime numbers and computing the composite  $n$  is performed, including for  $n$  that is more than 600 digits long, in less time than it takes to develop the pair of prime numbers  $p$  and  $q$  and compute that  $n$ .

32. (Amended) The method according to claim 18,

wherein  $p$  and  $q$  are a pair of prime numbers the product of which equals  $n$ ,

wherein the encoding step is divided into sub-steps, one sub-step for each of the  $k$  distinct random prime numbers,

wherein the  $k$  distinct random prime numbers are each smaller than  $p$  and  $q$ ,

whereby for a given length of  $n$  it takes fewer computational cycles to perform the encoding step relative to the number of computational cycles for performing such encoding step if the pair of prime numbers  $p$  and  $q$  were used instead.

33. (Amended) The method according to claim 18,

wherein  $p$  and  $q$  are a pair of prime numbers the product of which equals  $n$ , and

wherein developing the  $k$  distinct random prime numbers and computing the composite number  $n$  is performed, including for  $n$  that is more than 600 digits long, in less time than it takes to develop the pair of prime numbers  $p$  and  $q$  and compute that  $n$ .

34. (Amended) The method according to claim 14, wherein a message cryptographically processed by the sender with two-prime RSA public key encryption characterized by  $n$  being equal to a composite number computed as the product of 2 prime numbers  $p$  and  $q$ , is decipherable with multi-prime ( $k>2$ ) RSA public key encryption characterized by the composite number  $n$  being computed as the product of the  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ .

35. (Amended) The method according to claim 9, wherein the signed message word signal  $M_{ls}$ , formed from the digital message word signal  $M_l$  being cryptographically processed at the first terminal with multi-prime ( $k>2$ ) RSA public key encryption which is characterized by the composite number  $n$  being computed as the product of the  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ , is decipherable at the second terminal with two-prime RSA public key encryption characterized by  $n$  being equal to a composite number computed as the product of 2 prime numbers  $p$  and  $q$ .

36. (Amended) The method according to claim 16, wherein a message cryptographically processed by the sender with two-prime RSA public key encryption characterized by  $n$  being equal to a composite number computed as the product of 2 prime numbers  $p$  and  $q$ , is decipherable by the decoding with multi-prime ( $k>2$ ) RSA public key encryption characterized

by the composite number  $n$  being computed as the product of the  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ .

37. (Amended) The method according to claim 18, wherein the signed message  $M_s$ , formed from the plaintext message data  $M$  being cryptographically processed at the sender with multi-prime ( $k > 2$ ) RSA public key encryption which is characterized by the composite number  $n$  being computed as the product of the  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ , is decipherable by the decoding at the recipient with two-prime RSA public key encryption characterized by  $n$  being equal to a composite number computed as the product of 2 prime numbers  $p$  and  $q$ .

38. (Amended) The method according to claim 20, wherein a message cryptographically processed by the sender with two-prime RSA public key encryption characterized by  $n$  being equal to a composite number computed as the product of 2 prime numbers  $p$  and  $q$ , is decipherable at the recipient with multi-prime RSA public key encryption characterized by the composite number  $n$  being computed as the product of the at least three distinct random prime numbers.

39. (Amended) The method according to claim 22, wherein a message cryptographically processed by the sender with two-prime RSA public key encryption characterized by  $n$  being equal to a composite number computed as the product of 2 prime numbers  $p$  and  $q$ , is decipherable at the recipient with multi-prime RSA public key encryption characterized by the composite number  $n$  being computed as the product of the at least three distinct random prime numbers.

40. (Amended) A cryptography method for local storage of data by a private key owner, comprising the steps of:  
selecting a public key portion  $e$ ;  
developing  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ , where  $k \geq 3$ , and checking that each of the  $k$  distinct random prime numbers minus 1,  $p_1-1, p_2-1, \dots, p_k-1$ , is relatively prime to the public key portion  $e$ ;  
establishing a private key portion  $d$  by a relationship to the public key portion  $e$  in the form of

$$d \equiv e^{-1} \pmod{((p_1-1) \cdot (p_2-1) \cdots (p_k-1))};$$

computing a composite number,  $n$ , as a product of the  $k$  distinct random prime numbers that are factors of  $n$ , where only the private key owner knows the factors of  $n$ ; and encoding plaintext data  $M$  to ciphertext data  $C$  for the local storage, using a relationship of the form  $C \equiv M^e \pmod{n}$ , where  $0 \leq M \leq n-1$ , whereby the ciphertext data  $C$  is decipherable only by the private key owner having available to it the factors of  $n$ .

41. The cryptography method in accordance with claim 40, further comprising the step of: decoding the ciphertext data  $C$  from the local storage to the plaintext data  $M$  using a relationship of the form  $M \equiv C^d \pmod{n}$ .

42. (Amended) A cryptographic communications system, comprising:  
a plurality of stations;  
a communications medium; and  
a host system adapted to communicate with the plurality of stations via the communications medium sending a receiving messages cryptographically processed with an RSA public key encryption, the host system including  
at least one cryptosystem configured for  
developing  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ , where  $k \geq 3$ ,  
checking that each of the  $k$  distinct random prime numbers minus 1,  $p_1-1, p_2-1, \dots, p_k-1$ , is relatively prime to a public key portion  $e$  that is associated with the host system,  
computing a composite number,  $n$ , as a product of the  $k$  distinct random prime numbers,  
establishing a private key portion  $d$  by a relationship to the public key portion  $e$   
in the form of  $d \equiv e^{-1} \pmod{((p_1-1) \cdot (p_2-1) \cdots (p_k-1))}$ ,  
in response to an encoding request from the host system, encoding a plaintext message data  $M$  producing therefrom a ciphertext message data  $C$  to be communicated via the host system, the encoding using a relationship of the form  $C \equiv M^e \pmod{n}$ , where  $0 \leq M \leq n-1$ ,

in response to a decoding request from the host system, decoding a ciphertext message data  $C$  communicated via the host producing therefrom a plaintext message data  $M$  using a relationship of the form  $M \equiv C^d \pmod{n}$ .

43. (Amended) A system for communications of a message cryptographically processed with RSA public key encryption, comprising:

a bus; and

a cryptosystem communicatively coupled to and receiving from the bus encoding and decoding requests, the cryptosystem being configured for providing a public key portion  $e$ ,

developing  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ , where  $k \geq 3$ ,

checking that each of the  $k$  distinct random prime numbers minus 1,  $p_1-1, p_2-1, \dots, p_k-1$ , is relatively prime to the public key portion  $e$ ,

computing a composite number,  $n$ , as a product of the  $k$  distinct random prime numbers, establishing a private key portion  $d$  by a relationship to the public key portion  $e$  in the

form of  $d \equiv e^{-1} \pmod{(p_1-1) \cdot (p_2-1) \cdots (p_k-1)}$ ,

in response to an encoding request from the bus, encoding a plaintext form of a first message  $M$  to produce  $C$ , a ciphertext form of the first message, using a relationship of the form  $C \equiv M^e \pmod{n}$ , where  $0 \leq M \leq n-1$ , and

in response to an decoding request from the host system, decoding  $C'$ , a ciphertext form of a second message, to produce  $M'$ , a plaintext form of the second message, using a relationship of the form  $M' \equiv C'^d \pmod{n}$ , the first and second messages being distinct or one and the same.

44. The system of claim 42, wherein the at least one cryptosystem includes a plurality of exponentiators configured to operate in parallel in developing respective subtask values corresponding to the message.

45. (Amended) The system of claim 42, wherein the at least one cryptosystem includes a processor, a data-address bus,

a memory coupled to the processor via the data-address bus,  
a data encryption standard (DES) unit coupled the memory and the processor via the  
data-address bus,  
a plurality of exponentiator elements coupled to the processor via the DES unit, the  
plurality of exponentiator elements being configured to operate in parallel in  
developing respective subtask values corresponding to the message.

46. (Amended) The system of claim 45, wherein the memory and each of the plurality of exponentiator elements has its own DES unit that cryptographically processes message data received/returned from/to the processor.

47. (Amended) The system of claim 45, wherein the memory is partitioned into address spaces addressable by the processor, including secure, insecure and exponentiator elements address spaces, and wherein the DES unit is configured to recognize the secure and exponentiator elements address spaces and to automatically encode message data therefrom before it is provided to the exponentiator elements, the DES unit being bypassed when the processor is accessing the insecure memory address spaces, the DES unit being further configured to decode encoded message data received from the memory before it is provided to the processor.

48. The system of claim 45, wherein the at least one cryptosystem meets FIPS (Federal Information Processing Standard) 140-1 level 3.

49. The system of claim 45, wherein the processor maintains in the memory the public key portion  $e$  and the composite number  $n$  with its factors  $p_1, p_2, \dots, p_k$ .

50. (Amended) A system for communications of a message cryptographically processed with RSA public key encryption, comprising:  
a bus; and  
a cryptosystem receiving from the system via the bus encoding and decoding requests, the cryptosystem including  
a plurality of exponentiator elements configured to develop subtask values,

a memory, and

a processor configured for

receiving the encoding and decoding requests, each encoding request providing a plaintext message  $M$  to be encoded,

obtaining a public key that includes an exponent  $e$  and a modulus  $n$ , a representation of the modulus  $n$  existing in the memory in the form of its  $k$  distinct random prime number factors  $p_1, p_2, \dots, p_k$ , where  $k \geq 3$ ,

constructing subtasks, one subtask for each of the  $k$  factors, to be executed by the exponentiator elements for producing respective ones of the subtask values,  $C_1, C_2, \dots, C_k$ , and

forming a ciphertext message  $C$  from the subtask values  $C_1, C_2, \dots, C_k$ ,

wherein the ciphertext message  $C$  is decipherable using a private key that includes the modulus  $n$  and an exponent  $d$  which is a function of  $e$ .

51. (Amended) The system of claim 50 wherein each one of the subtasks  $C_1, C_2, \dots, C_k$  is developed using a relationship of the form  $C_i \equiv M_i^{e_i} \pmod{p_i}$ , where  $M_i \equiv M \pmod{p_i}$ , and  $e_i \equiv e \pmod{p_i - 1}$ , and where  $i=1, 2, \dots, k$ .

52. (Amended) A system for communications of a message cryptographically processed with RSA public key encryption, comprising:

a bus; and

a cryptosystem receiving from the system via the bus encoding and decoding requests, the cryptosystem including

a plurality of exponentiator elements configured to develop subtask values,

a memory, and

a processor configured for

receiving the encoding and decoding requests, each encoding/decoding request provided with a plaintext/ciphertext message  $M/C$  to be encoded/decoded and with or without a public/private key that includes an exponent  $e/d$  and a modulus  $n$  a representation of which exists in the memory in the form of its  $k$  distinct random prime number factors  $p_1, p_2, \dots, p_k$ , where  $k \geq 3$ ,

obtaining the public/private key from the memory if the encoding/decoding request is provided without the public/private key,  
constructing subtasks to be executed by the exponentiator elements for producing respective ones of the subtask values,  $M_1, M_2, \dots, M_k/C_1, C_2, \dots, C_k$ , and forming the ciphertext/plaintext message  $C/M$  from the subtask values  $C_1, C_2, \dots, C_k/M_1, M_2, \dots, M_k$ .

53. (Amended) The system of claim 52 wherein when produced each one of the subtasks  $C_1, C_2, \dots, C_k$  is developed using a relationship of the form  $C_i \equiv M_i^{e_i} \pmod{p_i}$ , where  $C_i \equiv C \pmod{p_i}$ , and  $e_i \equiv e \pmod{p_i - 1}$ , and where  $i=1, 2, \dots, k$ .

54. (Amended) The system of claim 52 wherein when produced each one of the subtasks  $M_1, M_2, \dots, M_k$  is developed using a relationship of the form  $M_i \equiv C_i^{d_i} \pmod{p_i}$ , where  $M_i \equiv M \pmod{p_i}$ , and  $d_i \equiv d \pmod{p_i - 1}$ , and where  $i=1, 2, \dots, k$ .

55. The system of claim 54, wherein the private key exponent  $d$  relates to the public key exponent  $e$  via  $d \equiv e^{-1} \pmod{((p_1 - 1) \cdot (p_2 - 1) \cdots (p_k - 1))}$ .

56. (Amended) A system for communications of a message cryptographically processed with RSA public key encryption, comprising:  
means for selecting a public key portion  $e$ ;  
means for developing  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ , where  $k \geq 3$ , and for checking that each of the  $k$  distinct random prime numbers minus 1,  $p_1-1, p_2-1, \dots, p_k-1$ , is relatively prime to the public key portion  $e$ ;  
means for establishing a private key portion  $d$  by a relationship to the public key portion  $e$  in the form of  $d \equiv e^{-1} \pmod{((p_1 - 1) \cdot (p_2 - 1) \cdots (p_k - 1))}$ ;  
means for computing a composite number,  $n$ , as a product of the  $k$  distinct random prime numbers;  
means for receiving a ciphertext message data  $C$ ; and

means for decoding the ciphertext message data  $C$  to a plaintext message data  $M$  using a relationship of the form  $M \equiv C^d \pmod{n}$ .

57. The system according to claim 56, further comprising:

means for encoding the plaintext message data  $M$  to the ciphertext message data  $C$ , using a relationship of the form  $C \equiv M^e \pmod{n}$ , where  $0 \leq M \leq n-1$ .

58. (Amended) A system for communications of a message cryptographically processed with RSA public key encryption, comprising:

means for selecting a public key portion  $e$ ;

means for developing  $k$  distinct random prime numbers,  $p_1, p_2, \dots, p_k$ , where  $k \geq 3$ , and for checking that each of the  $k$  distinct random prime numbers minus 1,  $p_1-1, p_2-1, \dots, p_k-1$ , is relatively prime to the public key portion  $e$ ;

means for establishing a private key portion  $d$  by a relationship to the public key portion  $e$  of the form  $d \equiv e^{-1} \pmod{((p_1-1) \cdot (p_2-1) \cdots (p_k-1))}$ ;

means for computing a composite number,  $n$ , as a product of the  $k$  distinct random prime numbers; and

means for encoding a plaintext message data  $M$  with the private key portion  $d$  to produce a signed message  $M_s$  using a relationship of the form  $M_s \equiv M^d \pmod{n}$ , where  $0 \leq M \leq n-1$ , the signed message  $M_s$  being decipherable using the public key portion  $e$ .

59. (Amended) The system of claim 58 further comprising the step of:

means for decoding the signed message  $M_s$  with the public key portion  $e$  to produce the plaintext message data  $M$  using a relationship of the form  $M \equiv M_s^e \pmod{n}$ .

60. (Amended) The system of claim 57, wherein the system can communicate the cryptographically processed message to another system that encodes/decodes data with RSA public key encryption using a modulus value equal to  $n$  independent of the  $k$  distinct prime numbers.

61. (Amended) The system of claim 59, wherein the system can communicate the cryptographically processed message to another system that encodes/decodes data with RSA public key encryption using a modulus value equal to  $n$  independent of the  $k$  distinct prime numbers.